UNIVERSITY OF CALIFORNIA, DAVIS

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EEC173B/ECS152C

Spring 2009

Homework 3 (due May 5)

1. (28 points) Mobile IP

- (a) (8 pt) List the entities of mobile IP and describe data transfer from a mobile node to a fixed node and vice versa. Why and where is encapsulation needed?
- (b) (5 pt) Show the steps required for a handover from one foreign agent to another foreign agent, including layer 2 and layer 3.
- (c) (5 pt) Explain packet flow if two mobile nodes communicate and both are in foreign networks. What additional routes do packets take if reverse tunneling is required?
- (d) (5 pt) Name the inefficiencies of mobile IP regarding data forwarding from a corresponding node to a mobile node. What optimizations are possible and what additional problems do they cause?
- (e) (5 pt) What is the basic purpose of DHCP and how can it be used for supporting mobility?

2. (22 points) Investigation of WLAN throughput.

- a) (5 pt) Consider an IEEE 802.11b network with nominal 11Mbps data rate. Assume you are the only user in the system that runs a specific application. Consider two cases:
 - A single TCP session is used to transmit the application data,
 - Four simultaneous TCP sessions are used to transmit the application data. Which approach achieves better throughput, and why?
- b) (7 pt) How and why does I-TCP isolate problems on the wireless link? What are the main drawbacks of this solution?
- c) Imagine a TCP session over wireless where the congestion window is fixed at 5 packets (congestion control is turned off and no fast retransmits). The receiver has infinite buffer and it sends an acknowledgment as soon as it receives a packet, i.e., acknowledgments are not deferred. Similarly sender transmits a packet as soon as it is allowed to. Each packet carries 1000 bytes and the time to transmit a packet is 2 ms. Assume that transmission of ACK takes negligible time. Note that the retransmission timer for a packet is started after the last bit of the packet is sent.

TCP's retransmission scheme is similar to go-back-N protocol, but it uses cumulative acknowledgements, i.e., when a receiver sends an ACK with sequence number N, it implies that all bytes from 0 to N-1 have been received correctly.

- i. (5 pt) Suppose two *data* segments with byte sequence numbers 3000 and 15000 are lost once during the transmission. How many segments get retransmitted under each of the following conditions?
 - Round trip time = 100 ms, Timeout = 101 ms
 - Round trip time = 100 ms, Timeout = 152 ms
- ii. (5 pt) Suppose *acknowledgments* corresponding to the above data segments are lost instead of the data segments. How many segments get retransmitted under the above conditions?
 - Round trip time = 100 ms, Timeout = 101 ms
 - Round trip time = 100 ms, Timeout = 152 ms