Instruction Level Parallelism and Its Dynamic Exploitation

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Based on Orginal Lecture Material from
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How to Improve CPI?

a) In-order issue and In-order completion is rather limited
b) Allow out-of-order issue
c) Allow out-of-order completion
d) Allow both
e) Multiple Issue
f) Basic block is small – 4-7 instructions before we hit a control flow instruction
g) Look for ILP across basic blocks
h) Predict Branches
i) Speculation

Exploit instruction level parallelism
Have multiple functional units to execute more than one instruction at the same time

Basic Problems to address

- Detecting data hazards – with larger window sizes becomes more complex
- Deal with control hazards
- Deal with out-of-order execution
- Deal with out-of-order completion
- Precise Interrupts

Two Philosophies

- Static (Let the compiler figure these out and deal with them) – Chapter 4
- Dynamic (Let the hardware figure it out) – Chapter 3
- Both, yes, the familiar co-design again – Itanium approach

Roadmap

	Technique	Reduces
Chapter 3	Dynamic scheduling	Data hazard stalls
	Dynamic branch prediction	Control stalls
	Issuing multiple instructions per cycle	Ideal CPI
	Speculation	Data and control stalls
	Dynamic memory disambiguation	Data hazard stalls involving memory
	Loop unrolling	Control hazard stalls
Chapter 4	Basic compiler pipeline scheduling	Data hazard stalls
	Compiler dependence analysis	Ideal CPI and data hazard stalls
	Software pipelining and trace scheduling	Ideal CPI and data hazard stalls
	Compiler speculation	Ideal CPI, data and control stalls

Instruction-Level Parallelism (ILP)

- Basic Block (BB) ILP is quite small
  - BB: a straight-line code sequence with no branches in, except to the entry and no branches out except at the exit
  - average dynamic branch frequency 15% to 25%
    => 4 to 7 instructions execute between a pair of branches
  - Plus instructions in BB likely to depend on each other
- To obtain substantial performance enhancements, we must exploit ILP across multiple basic blocks
- Simplest: loop-level parallelism to exploit parallelism among iterations of a loop
  - Vector is one way
  - If not vector, then either dynamic via branch prediction or static via loop unrolling by compiler

Data Dependence and Hazards

- $Instr_J$ is data dependent on $Instr_I$
  $Instr_J$ tries to read operand before $Instr_I$ writes it

```
I: add r1,r2,r3
J: sub r4,r1,r3
```

- or $Instr_J$ is data dependent on $Instr_K$ which is dependent on $Instr_I$
- Caused by a "True Dependence" (compiler term)
- If true dependence causes a hazard in the pipeline, it is called a Read After Write (RAW) hazard

Data Dependence and Hazards

- Dependences are a property of programs
- Presence of dependence indicates potential for a hazard, but actual hazard and length of any stall is a property of the pipeline
- Importance of the data dependencies

1) indicates the possibility of a hazard
2) determines order in which results must be calculated
3) sets an upper bound on how much parallelism can possibly be exploited

Name Dependence #1: Anti-dependence

- Name dependence: when 2 instructions use same register or memory location, called a name, but no flow of data between the instructions associated with that name; 2 versions of name dependence
- $Instr_J$ writes operand *before* $Instr_I$ reads it

```
I: sub r4,r1,r3
J: add r1,r2,r3
K: mul r6,r1,r7
```

- Called an "anti-dependence" in compiler work.
  This results from reuse of the name "r1"
- If anti-dependence caused a hazard in the pipeline, called a Write After Read (WAR) hazard

Name Dependence #2: Output dependence

- $Instr_J$ writes operand *before* $Instr_I$ writes it.

```
I: sub r1,r4,r3
J: add r1,r2,r3
K: mul r6,r1,r7
```

- Called an "output dependence" by compiler writers This also results from the reuse of name "r1"
- If anti-dependence caused a hazard in the pipeline, called a Write After Write (WAW) hazard

Memory Dataflow

- When instructions communicate through memory
- Memory aliasing is a problem
- Eg: store val, 100(r4)
  load r5, 20(r6)

How do you know if 100(r4) and 20(r6) are the pointing to the same physical memory location or not?

Impossible to determine statically if two pointers are pointing to the same physical memory location

Dynamic memory disambiguation is needed which adds overhead – runtime check.

ILP and Data Hazards

- program order: order instructions would execute in if executed sequentially 1 at a time as determined by original source program
- HW/SW goal: exploit parallelism by preserving appearance of program order
  - modify order in manner than cannot be observed by program
  - must not affect the outcome of the program
- Ex: Instructions involved in a name dependence can execute simultaneously if name used in instructions is changed so instructions do not conflict
  - Register renaming resolves name dependence for registers
  - Either by compiler or by HW
  - add r1, r2, r3
  - sub r2, r4,r5
  - and r3, r2, 1

Control Dependencies

- Every instruction is control dependent on some set of branches, and, in general, these control dependencies must be preserved to preserve program order

```
if p1 {
  S1;
};
if p2 {
  S2;
}
```

- S1 is control dependent on p1, and S2 is control dependent on p2 but not on p1.

Control Dependence CAN Ignored

- Control dependence need not always be preserved
  - willing to execute instructions that should not have been executed, thereby violating the control dependences,

    if can do so without affecting correctness of the program
- 2 properties critical to program correctness are exception behavior and data flow

Exception Behavior

- Preserving exception behavior => any changes in instruction execution order must not change how exceptions are raised in program (=> no new exceptions)
- Example:

```
    DADDU     R2,R3,R4
    BEQZ      R2,L1
    LW        R1,0(R2)
L1:
```

- Problem with moving LW before BEQZ?

Data Flow

- Data flow: actual flow of data values among instructions that produce results and those that consume them
  - branches make flow dynamic, determine which instruction is supplier of data
- Example:

```
DADDU     R1,R2,R3
BEQZ      R4,L
DSUBU     R1,R5,R6
L:  …
OR        R7,R1,R8
```

- OR depends on DADDU or DSUBU?
  Must preserve data flow on execution

Dynamic Scheduling

- Handles cases when dependences unknown at compile time
  - (e.g., because they may involve a memory reference)
- It simplifies the compiler
- Allows code that compiled for one pipeline to run efficiently on a different pipeline
- Hardware speculation, a technique with significant performance advantages, that builds on dynamic scheduling

HW Schemes: Instruction Parallelism

- Key idea: Allow instructions behind stall to proceed

```
DIVD  F0,F2,F4
ADDD  F10,F0,F8
SUBD  F12,F8,F14
```

- Enables out-of-order execution and allows out-of-order completion
- Will distinguish when an instruction *begins execution* and when it *completes execution*; between those two times, the instruction is *in execution*
- In a dynamically scheduled pipeline, all instructions pass through issue stage in order (in-order issue)

Dynamic Scheduling Step 1

- Simple pipeline had 1 stage to check both structural and data hazards: Instruction Decode (ID), also called Instruction Issue
- Split the ID pipe stage of simple 5-stage pipeline into 2 stages:
- *Issue*—Decode instructions, check for structural hazards
- *Read operands*—Wait until no data hazards, then read operands

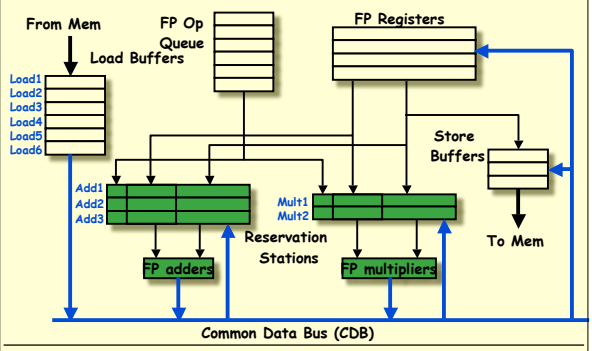
A Dynamic Algorithm: Tomasulo's Algorithm

- For IBM 360/91 (before caches!)
- Goal: High Performance without special compilers
- Small number of floating point registers (4 in 360) prevented interesting compiler scheduling of operations
  - This led Tomasulo to try to figure out how to get more effective registers — renaming in hardware!
- Why Study 1966 Computer?
- The descendants of this have flourished!
  - Alpha 21264, HP 8000, MIPS 10000, Pentium III, PowerPC 604, …

Tomasulo Algorithm

- Control & buffers distributed with Function Units (FU)
  - FU buffers called "reservation stations"; have pending operands
- Registers in instructions replaced by values or pointers to reservation stations(RS);
  - form of register renaming ;
  - avoids WAR, WAW hazards
  - More reservation stations than registers, so can do optimizations compilers can't
- Results to FU from RS, not through registers, over Common Data Bus that broadcasts results to all FUs
- Load and Stores treated as FUs with RSs as well
- Integer instructions can go past branches, allowing FP ops beyond basic block in FP queue

Tomasulo Organization



Reservation Station Components

Op: Operation to perform in the unit (e.g., + or –)

Vj, Vk: Value of Source operands

- Store buffers has V field, result to be stored

Qj, Qk: Reservation stations producing source registers (value to be written)

- Note: Qj,Qk=0 => ready
- Store buffers only have Qi for RS producing result

Busy: Indicates reservation station or FU is busy

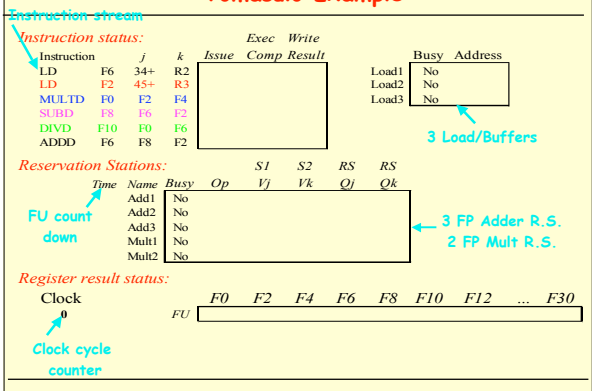
Register result status—Indicates which functional unit will write each register, if one exists. Blank when no pending instructions that will write that register.

Three Stages of Tomasulo Algorithm

1. Issue—get instruction from FP Op Queue
   If reservation station free (no structural hazard), control issues instr & sends operands (renames registers).
2. Execute—operate on operands (EX)
   When both operands ready then execute; if not ready, watch Common Data Bus for result
3. Write result—finish execution (WB)
   Write on Common Data Bus to all awaiting units; mark reservation station available

- Normal data bus: data + destination ("go to" bus)
- Common data bus: data + source ("come from" bus)
  - 64 bits of data + 4 bits of Functional Unit source address
  - Write if matches expected Functional Unit (produces result)
  - Does the broadcast
- Example speed:
  3 clocks for Fl .pt. +,-; 10 for * ; 40 clks for /

Tomasulo Example



Tomasulo Example Cycle 1

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1		
LD F2	45+	R3			
MULTD F0	F2	F4			
SUBD F8	F6	F2			
DIVD F10	F0	F6			
ADDD F6	F8	F2			

	Busy	Address
Load1	Yes	34+R2
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
	Mult2	No					

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
1	FU			Load1					

Tomasulo Example Cycle 2

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1		
LD F2	45+	R3	2		
MULTD F0	F2	F4			
SUBD F8	F6	F2			
DIVD F10	F0	F6			
ADDD F6	F8	F2			

	Busy	Address
Load1	Yes	34+R2
Load2	Yes	45+R3
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
	Mult2	No					

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
2	FU	Load2		Load1					

Note: Can have multiple loads outstanding

Tomasulo Example Cycle 3

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	
LD F2	45+	R3	2		
MULTD F0	F2	F4	3		
SUBD F8	F6	F2			
DIVD F10	F0	F6			
ADDD F6	F8	F2			

	Busy	Address
Load1	Yes	34+R2
Load2	Yes	45+R3
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	MULTD		R(F4)		Load2
	Mult2	No					

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
3	FU	Mult1	Load2		Load1				

- Note: registers names are removed ("renamed") in Reservation Stations; MULT issued
- ~~Load1 completing; what is waiting for Load1?~~

Tomasulo Example Cycle 4

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4		
DIVD F10	F0	F6			
ADDD F6	F8	F2			

	Busy	Address
Load1	No	
Load2	Yes	45+R3
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	Yes	SUBD	M(A1)			Load2
	Add2	No					
	Add3	No					
	Mult1	Yes	MULTD		R(F4)		Load2
	Mult2	No					

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
4	FU	Mult1	Load2		M(A1)	Add1			

- Load2 completing; what is waiting for Load2?

Tomasulo Example Cycle 5

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4		
DIVD F10	F0	F6	5		
ADDD F6	F8	F2			

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
2	Add1	Yes	SUBD	M(A1)	M(A2)		
	Add2	No					
	Add3	No					
10	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
5	FU Mult1	M(A2)		M(A1)	Add1	Mult2			

- Timer starts down for Add1, Mult1

Tomasulo Example Cycle 6

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4		
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6		

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
1	Add1	Yes	SUBD	M(A1)	M(A2)		
	Add2	Yes	ADDD		M(A2)	Add1	
	Add3	No					
9	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
6	FU Mult1	M(A2)		Add2	Add1	Mult2			

- Issue ADDD here despite name dependency on F6?

Tomasulo Example Cycle 7

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6		

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
0	Add1	Yes	SUBD	M(A1)	M(A2)		
	Add2	Yes	ADDD		M(A2)	Add1	
	Add3	No					
8	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
7	FU Mult1	M(A2)		Add2	Add1	Mult2			

- Add1 (SUBD) completing; what is waiting for it?

Tomasulo Example Cycle 8

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6		

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
2	Add2	Yes	ADDD	M-M1	M(A2)		
	Add3	No					
7	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
8	FU Mult1	M(A2)		Add2	M-M1	Mult2			

Tomasulo Example Cycle 9

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6		

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
1	Add2	Yes	ADDD	(M-M)	M(A2)		
	Add3	No					
6	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
9	FU Mult1	M(A2)		Add2	(M-M)	Mult2			

Tomasulo Example Cycle 10

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6	10	

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
0	Add2	Yes	ADDD	(M-M)	M(A2)		
	Add3	No					
5	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
10	FU Mult1	M(A2)		Add2	(M-M)	Mult2			

- Add2 (ADDD) completing; what is waiting for it?

Tomasulo Example Cycle 11

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
4	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
11	FU Mult1	M(A2)		(M-M+M	(M-M)	Mult2			

- Write result of ADDD here?
- All quick instructions complete in this cycle!

Tomasulo Example Cycle 12

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
3	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
12	FU Mult1	M(A2)		(M-M+M	(M-M)	Mult2			

Tomasulo Example Cycle 13

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
2	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
13	FU Mult1	M(A2)		(M-M+M	(M-M)	Mult2			

Tomasulo Example Cycle 14

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3		
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
1	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
14	FU Mult1	M(A2)		(M-M+M	(M-M)	Mult2			

Tomasulo Example Cycle 15

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3	15	
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
0	Mult1	Yes	MULTD	M(A2)	R(F4)		
	Mult2	Yes	DIVD		M(A1)	Mult1	

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
15	FU Mult1	M(A2)		(M-M+M	(M-M)	Mult2			

- Mult1 (MULTD) completing; what is waiting for it?

Tomasulo Example Cycle 16

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD F6	34+	R2	1	3	4
LD F2	45+	R3	2	4	5
MULTD F0	F2	F4	3	15	16
SUBD F8	F6	F2	4	7	8
DIVD F10	F0	F6	5		
ADDD F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
40	Mult2	Yes	DIVD	M*F4	M(A1)		

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
16	FU M*F4	M(A2)		(M-M+M	(M-M)	Mult2			

- Just waiting for Mult2 (DIVD) to complete

Faster than light computation (skip a couple of cycles)

Tomasulo Example Cycle 55

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD	F6	34+	R2	1	3	4
LD	F2	45+	R3	2	4	5
MULTD	F0	F2	F4	3	15	16
SUBD	F8	F6	F2	4	7	8
DIVD	F10	F0	F6	5		
ADDD	F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
1	Mult2	Yes	DIVD	M*F4	M(A1)		

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
55	FU	M*F4	M(A2)		(M-M+M	(M-M)	Mult2		

Tomasulo Example Cycle 56

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD	F6	34+	R2	1	3	4
LD	F2	45+	R3	2	4	5
MULTD	F0	F2	F4	3	15	16
SUBD	F8	F6	F2	4	7	8
DIVD	F10	F0	F6	5	56	
ADDD	F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
0	Mult2	Yes	DIVD	M*F4	M(A1)		

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
56	FU	M*F4	M(A2)		(M-M+M	(M-M)	Mult2		

- Mult2 (DIVD) is completing; what is waiting for it?

Tomasulo Example Cycle 57

Instruction status:

Instruction	j	k	Issue	Exec Comp	Write Result
LD	F6	34+	R2	1	3	4
LD	F2	45+	R3	2	4	5
MULTD	F0	F2	F4	3	15	16
SUBD	F8	F6	F2	4	7	8
DIVD	F10	F0	F6	5	56	57
ADDD	F6	F8	F2	6	10	11

	Busy	Address
Load1	No	
Load2	No	
Load3	No	

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
	Mult2	Yes	DIVD	M*F4	M(A1)		

Register result status:

Clock	F0	F2	F4	F6	F8	F10	F12	...	F30
56	FU	M*F4	M(A2)		(M-M+M	(M-M)	Result		

- Once again: In-order issue, out-of-order execution and out-of-order completion.

Tomasulo Drawbacks

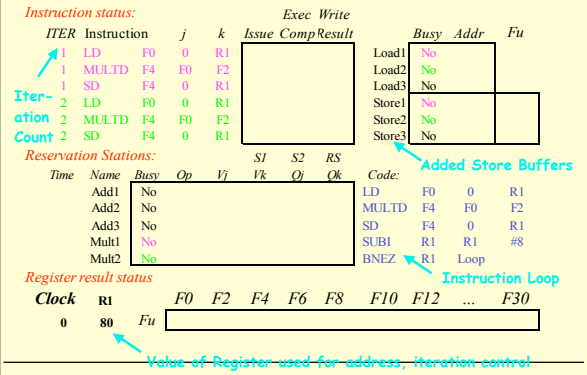
- Complexity
  - delays of 360/91, MIPS 10000, Alpha 21264, IBM PPC 620 in CA:AQA 2/e, but not in silicon!
- Many associative stores (CDB) at high speed
- Performance limited by Common Data Bus
  - Each CDB must go to multiple functional units ⇒high capacitance, high wiring density
  - Number of functional units that can complete per cycle limited to one!
    - » Multiple CDBs ⇒ more FU logic for parallel assoc stores
- Non-precise interrupts!
  - We will address this later

Tomasulo Loop Example

```
Loop:   LD      F0    0    R1
        MULTD   F4    F0   F2
        SD      F4    0    R1
        SUBI    R1    R1   #8
        BNEZ    R1    Loop
```

- This time assume Multiply takes 4 clocks
- Assume 1st load takes 8 clocks (L1 cache miss), 2nd load takes 1 clock (hit)
- To be clear, will show clocks for SUBI, BNEZ
  - Reality: integer instructions ahead of Fl. Pt. Instructions
- Show 2 iterations

Loop Example



Loop Example Cycle 1

Loop Example Cycle 2

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1		
1	MULTD F4	F0	F2	2		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	No		
Load3	No		
Store1	No		
Store2	No		
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	No					

Code:

LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
2	80	Fu	Load1		Mult1					

Loop Example Cycle 3

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1		
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	No		
Load3	No		
Store1	Yes	80	Mult1
Store2	No		
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	No					

Code:

LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
3	80	Fu	Load1		Mult1					

- Implicit renaming sets up data flow graph

Loop Example Cycle 4

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1		
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	No		
Load3	No		
Store1	Yes	80	Mult1
Store2	No		
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	No					

Code:

LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
4	80	Fu	Load1		Mult1					

- Dispatching SUBI Instruction (not in FP queue)

Loop Example Cycle 5

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1		
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	No		
Load3	No		
Store1	Yes	80	Mult1
Store2	No		
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	No					

Code:

LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
5	72	Fu	Load1		Mult1					

- And, BNEZ instruction (not in FP queue)

Loop Example Cycle 6

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1		
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	Yes	72	
Load3	No		
Store1	Yes	80	Mult1
Store2	No		
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	No					

Code:

```
LD    F0  0     R1
MULTD F4  F0    F2
SD    F4  0     R1
SUBI  R1  R1    #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
6	72	Load2		Mult1						

- Notice that F0 never sees Load from location 80

Loop Example Cycle 7

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1		
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6		
2	MULTD F4	F0	F2	7		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	Yes	72	
Load3	No		
Store1	Yes	80	Mult1
Store2	No		
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	Yes	Multd		R(F2)	Load2	

Code:

```
LD    F0  0     R1
MULTD F4  F0    F2
SD    F4  0     R1
SUBI  R1  R1    #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
7	72	Load2		Mult2						

- Register file completely detached from computation
- First and Second iteration completely overlapped

Loop Example Cycle 8

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1		
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6		
2	MULTD F4	F0	F2	7		
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	Yes	72	
Load3	No		
Store1	Yes	80	Mult1
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	Yes	Multd		R(F2)	Load2	

Code:

```
LD    F0  0     R1
MULTD F4  F0    F2
SD    F4  0     R1
SUBI  R1  R1    #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
8	72	Load2		Mult2						

Loop Example Cycle 9

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6		
2	MULTD F4	F0	F2	7		
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	Yes	80	
Load2	Yes	72	
Load3	No		
Store1	Yes	80	Mult1
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load1	
	Mult2	Yes	Multd		R(F2)	Load2	

Code:

```
LD    F0  0     R1
MULTD F4  F0    F2
SD    F4  0     R1
SUBI  R1  R1    #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
9	72	Load2		Mult2						

- Load1 completing: who is waiting?
- Note: Dispatching SUBI

Loop Example Cycle 10

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	
2	MULTD F4	F0	F2	7		
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	Yes	72	
Load3	No		
Store1	Yes	80	Mult1
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
4	Mult1	Yes	Multd	M[80]	R(F2)		
	Mult2	Yes	Multd		R(F2)	Load2	

Code:

```
LD    F0  0   R1
MULTD F4  F0  F2
SD    F4  0   R1
SUBI  R1  R1  #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
10	64	Load2		Mult2						

- Load2 completing: who is waiting?
- Note: Dispatching BNEZ

Loop Example Cycle 11

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7		
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	Mult1
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
3	Mult1	Yes	Multd	M[80]	R(F2)		
4	Mult2	Yes	Multd	M[72]	R(F2)		

Code:

```
LD    F0  0   R1
MULTD F4  F0  F2
SD    F4  0   R1
SUBI  R1  R1  #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
11	64	Load3		Mult2						

- Next load in sequence

Loop Example Cycle 12

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7		
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	Mult1
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
2	Mult1	Yes	Multd	M[80]	R(F2)		
3	Mult2	Yes	Multd	M[72]	R(F2)		

Code:

```
LD    F0  0   R1
MULTD F4  F0  F2
SD    F4  0   R1
SUBI  R1  R1  #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
12	64	Load3		Mult2						

- Why not issue third multiply?

Loop Example Cycle 13

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2		
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7		
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	Mult1
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	S1 Vj	S2 Vk	RS Qj	RS Qk
	Add1	No					
	Add2	No					
	Add3	No					
1	Mult1	Yes	Multd	M[80]	R(F2)		
2	Mult2	Yes	Multd	M[72]	R(F2)		

Code:

```
LD    F0  0   R1
MULTD F4  F0  F2
SD    F4  0   R1
SUBI  R1  R1  #8
BNEZ  R1  Loop
```

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
13	64	Load3		Mult2						

- Why not issue third store?

Loop Example Cycle 14

Instruction status:

ITER	Instruction	j	k	Issue	Comp	Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2	14	
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7		
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	Mult1
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	Vj	Vk	Qj	Qk
	Add1	No					
	Add2	No					
	Add3	No					
0	Mult1	Yes	Multd	M[80]	R(F2)		
1	Mult2	Yes	Multd	M[72]	R(F2)		

Code:			
LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
14	64	Fu	Load3		Mult2					

- Mult1 completing. Who is waiting?

Loop Example Cycle 15

Instruction status:

ITER	Instruction	j	k	Issue	Comp	Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2	14	15
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7	15	
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	80]*R2
Store2	Yes	72	Mult2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	Vj	Vk	Qj	Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	No					
0	Mult2	Yes	Multd	M[72]	R(F2)		

Code:			
LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
15	64	Fu	Load3		Mult2					

- Mult2 completing. Who is waiting?

Loop Example Cycle 16

Instruction status:

ITER	Instruction	j	k	Issue	Comp	Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2	14	15
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7	15	16
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	80]*R2
Store2	Yes	72	72]*R2
Store3	No		

Reservation Stations:

Time	Name	Busy	Op	Vj	Vk	Qj	Qk
	Add1	No					
	Add2	No					
	Add3	No					
4	Mult1	Yes	Multd		R(F2)	Load3	
	Mult2	No					

Code:			
LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
16	64	Fu	Load3		Mult1					

Loop Example Cycle 17

Instruction status:

ITER	Instruction	j	k	Issue	Comp	Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2	14	15
1	SD F4	0	R1	3		
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7	15	16
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	80]*R2
Store2	Yes	72	72]*R2
Store3	Yes	64	Mult1

Reservation Stations:

Time	Name	Busy	Op	Vj	Vk	Qj	Qk
	Add1	No					
	Add2	No					
	Add3	No					
	Mult1	Yes	Multd		R(F2)	Load3	
	Mult2	No					

Code:			
LD	F0	0	R1
MULTD	F4	F0	F2
SD	F4	0	R1
SUBI	R1	R1	#8
BNEZ	R1	Loop	

Register result status

Clock	R1	F0	F2	F4	F6	F8	F10	F12	...	F30
17	64	Fu	Load3		Mult1					

Loop Example Cycle 18

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2	14	15
1	SD F4	0	R1	3	18	
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7	15	16
2	SD F4	0	R1	8		

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	Yes	80	[80]*R2
Store2	Yes	72	[72]*R2
Store3	Yes	64	Mult1

Reservation Stations:

Time	Name	Busy	Op	Vj	S1 Vk	S2 Qj	RS Qk	Code:			
	Add1	No						LD	F0	0	R1
	Add2	No						MULTD	F4	F0	F2
	Add3	No						SD	F4	0	R1
	Mult1	Yes	Multd		R(F2)	Load3		SUBI	R1	R1	#8
	Mult2	No						BNEZ	R1	Loop	

Register result status

Clock	R1		F0	F2	F4	F6	F8	F10	F12	...	F30
18	64	Fu	Load3		Mult1						

Loop Example Cycle 19

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2	14	15
1	SD F4	0	R1	3	18	19
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7	15	16
2	SD F4	0	R1	8	19	

	Busy	Addr	Fu
Load1	No		
Load2	No		
Load3	Yes	64	
Store1	No		
Store2	Yes	72	[72]*R2
Store3	Yes	64	Mult1

Reservation Stations:

Time	Name	Busy	Op	Vj	S1 Vk	S2 Qj	RS Qk	Code:			
	Add1	No						LD	F0	0	R1
	Add2	No						MULTD	F4	F0	F2
	Add3	No						SD	F4	0	R1
	Mult1	Yes	Multd		R(F2)	Load3		SUBI	R1	R1	#8
	Mult2	No						BNEZ	R1	Loop	

Register result status

Clock	R1		F0	F2	F4	F6	F8	F10	F12	...	F30
19	56	Fu	Load3		Mult1						

Loop Example Cycle 20

Instruction status:

ITER	Instruction	j	k	Issue	Exec Comp	Write Result
1	LD F0	0	R1	1	9	10
1	MULTD F4	F0	F2	2	14	15
1	SD F4	0	R1	3	18	19
2	LD F0	0	R1	6	10	11
2	MULTD F4	F0	F2	7	15	16
2	SD F4	0	R1	8	19	20

	Busy	Addr	Fu
Load1	Yes	56	
Load2	No		
Load3	Yes	64	
Store1	No		
Store2	No		
Store3	Yes	64	Mult1

Reservation Stations:

Time	Name	Busy	Op	Vj	S1 Vk	S2 Qj	RS Qk	Code:			
	Add1	No						LD	F0	0	R1
	Add2	No						MULTD	F4	F0	F2
	Add3	No						SD	F4	0	R1
	Mult1	Yes	Multd		R(F2)	Load3		SUBI	R1	R1	#8
	Mult2	No						BNEZ	R1	Loop	

Register result status

Clock	R1		F0	F2	F4	F6	F8	F10	F12	...	F30
20	56	Fu	Load1		Mult1						

- Once again: In-order issue, out-of-order execution and out-of-order completion.

Why can Tomasulo overlap iterations of loops?

- **Register renaming**
  - Multiple iterations use different physical destinations for registers (dynamic loop unrolling).
- **Reservation stations**
  - Permit instruction issue to advance past integer control flow operations
  - Also buffer old values of registers - totally avoiding the WAR stall that we saw in the scoreboard.
- **Other perspective: Tomasulo building data flow dependency graph on the fly.**

Tomasulo's scheme offers 2 major advantages

(1) the distribution of the hazard detection logic
- distributed reservation stations and the CDB
- If multiple instructions waiting on single result, & each instruction has other operand, then instructions can be released simultaneously by broadcast on CDB
- If a centralized register file were used, the units would have to read their results from the registers when register buses are available.

(2) the elimination of stalls for WAW and WAR hazards

What about Precise Interrupts?

- State of machine looks as if no instruction beyond faulting instructions has issued
- Tomasulo had:

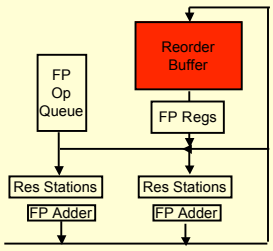
  In-order issue, out-of-order execution, and out-of-order completion
- Need to "fix" the out-of-order completion aspect so that we can find precise breakpoint in instruction stream.

Relationship between precise interrupts and speculation

- Speculation: guess and check
- Important for branch prediction:
  - Need to "take our best shot" at predicting branch direction.
- If we speculate and are wrong, need to back up and restart execution to point at which we predicted incorrectly:
  - This is exactly same as precise exceptions!
- Technique for both precise interrupts/exceptions and speculation: *in-order completion or commit*

HW support for precise interrupts

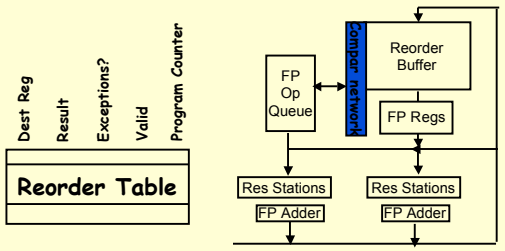
- Need HW buffer for results of uncommitted instructions: *reorder buffer*
  - 3 fields: instr, destination, value
  - Use reorder buffer number instead of reservation station when execution completes
  - Supplies operands between execution complete & commit
  - (Reorder buffer can be operand source => more registers like RS)
  - Instructions commit
  - Once instruction commits, result is put into register
  - As a result, easy to undo speculated instructions on mispredicted branches or exceptions



Four Steps of Speculative Tomasulo Algorithm

1. Issue—get instruction from FP Op Queue
   If reservation station and reorder buffer slot free, issue instr & send operands & reorder buffer no. for destination (this stage sometimes called "dispatch")
2. Execution—operate on operands (EX)
   When both operands ready then execute; if not ready, watch CDB for result; when both in reservation station, execute; checks RAW (sometimes called "issue")
3. Write result—finish execution (WB)
   Write on Common Data Bus to all awaiting FUs & reorder buffer; mark reservation station available.
4. Commit—update register with reorder result
   When instr. at head of reorder buffer & result present, update register with result (or store to memory) and remove instr from reorder buffer. Mispredicted branch flushes reorder buffer (sometimes called "graduation")

What are the hardware complexities with reorder buffer (ROB)?



- How do you find the latest version of a register?
  - (As specified by Smith paper) need associative comparison network
  - Could use future file or just use the register result status buffer to track which specific reorder buffer has received the value
- Need as many ports on ROB as register file

Summary

- Reservations stations: *implicit register renaming* to larger set of registers + buffering source operands
  - Prevents registers as bottleneck
  - Avoids WAR, WAW hazards of Scoreboard
  - Allows loop unrolling in HW
- Not limited to basic blocks (integer units gets ahead, beyond branches)
- Today, helps cache misses as well
  - Don't stall for L1 Data cache miss (insufficient ILP for L2 miss?)
- Lasting Contributions
  - Dynamic scheduling
  - Register renaming
  - Load/store disambiguation
- 360/91 descendants are Pentium III; PowerPC 604; MIPS R10000; HP-PA 8000; Alpha 21264