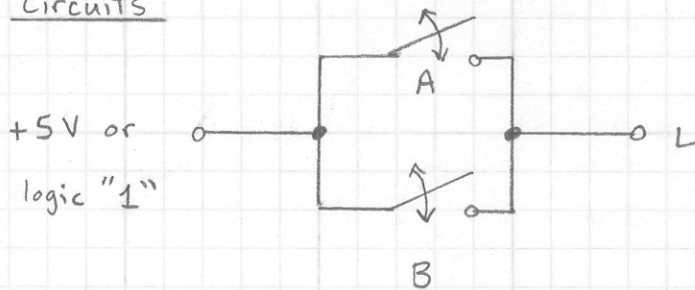


I. Digital Logic

Consider circuit elements in digital logic circuits as ideal switches (typically implemented using transistors.)

Circuits

Parallel Connection

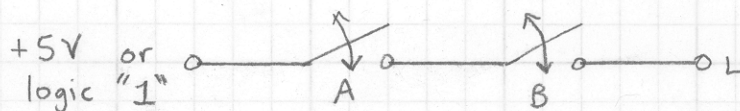
- When A, B are "1", then switches are closed.
- When A, B are "0", then switches are open.

Truth Table (mapping from input signals to output signal)

A	B	L
0	0	0
0	1	1
1	0	1
1	1	1

Mathematical (Logical) "OR" Operation

$$L = A + B$$

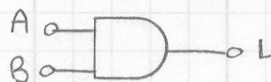


Series Connection

Truth Table

A	B	L
0	0	0
0	1	0
1	0	0
1	1	1

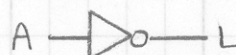
"AND" Operation



Mathematical Form
 $A \cdot B = L$ or $AB = L$

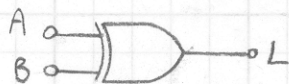
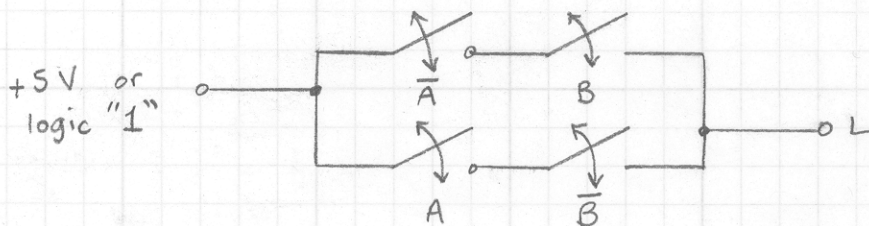
"NOT" Operation (Inversion)

A	L
0	1
1	0



Mathematical Form
 $L = \bar{A}$

"XOR" (Exclusive OR) Operation



Mathematical Form

$$L = (\bar{A} + \bar{B})(A + B) = \bar{A} \cdot B + A \cdot \bar{B}$$

$$L = A \oplus B$$

II. Boolean Algebra

Mathematical manipulations of equations involving single-bit binary numbers.

Basic Operations

"OR" +

"AND" •

"NOT" \bar{A}

Commutative Law

$$A + B = B + A \quad AB = BA$$

Associative Law

$$(A + B) + C = A + (B + C) \quad (AB)C = A(BC)$$

Distributive Law

$$A(B + C) = AB + AC$$

$$A + BC = (A + B)(A + C)$$

Identities:

$$A \cdot 1 = A$$

$$A + 0 = A$$

$$\bar{\bar{A}} = A$$

$$A \cdot 0 = 0$$

$$A + A = A$$

$$A \cdot A = A$$

$$A + 1 = 1$$

$$A \cdot \bar{A} = 0$$

$$A + \bar{A} = 1$$

Theorems of Boolean Algebra (pp. 772-778 in book)

$$A + AB = A$$

$$A(A+B) = A$$

$$A + \bar{A}B = A + B$$

$$A(\bar{A} + B) = AB$$

DeMorgan's Theorem

$$\overline{A+B} = \bar{A}\bar{B}$$

$$\overline{AB} = \bar{A} + \bar{B}$$

Sum of Products

Express Boolean function as OR of ANDs

$$F = \bar{A}\bar{B}C + \bar{A}B\bar{C} + A\bar{B}\bar{C} + \underbrace{ABC}_{\text{minterm}}$$

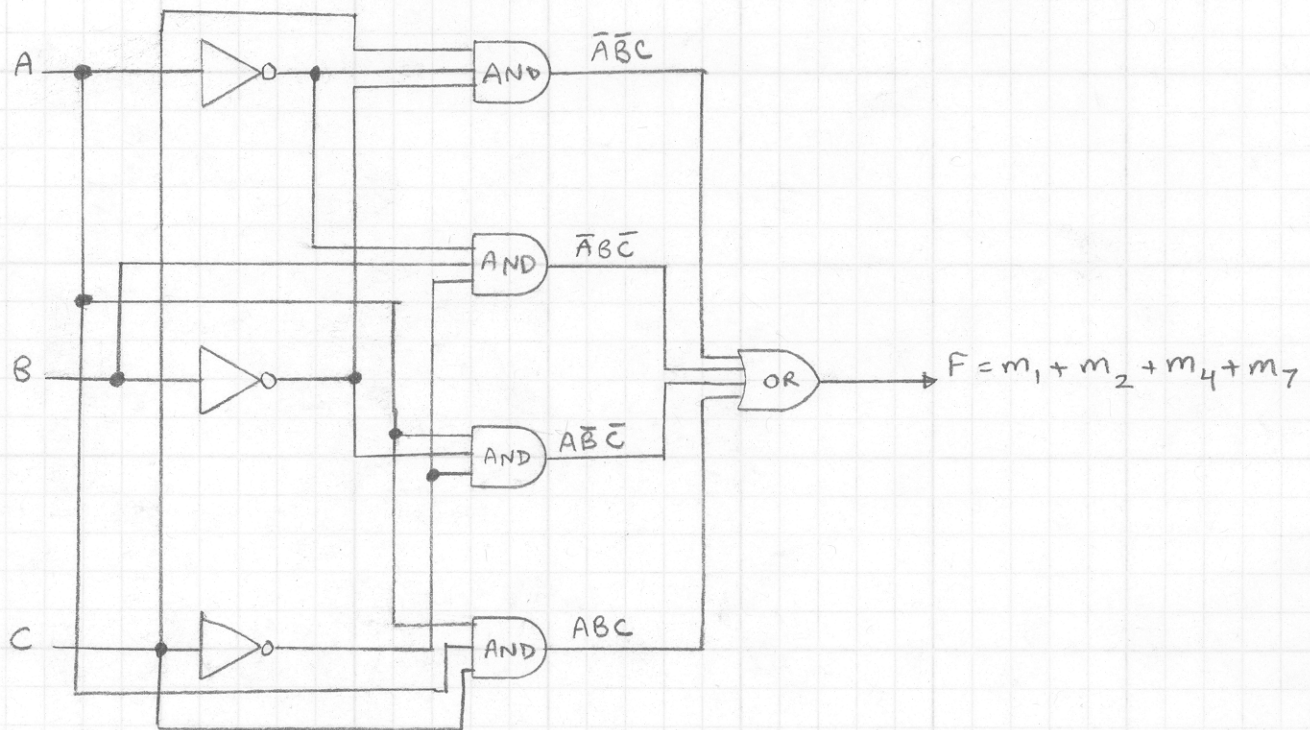
Think about
where $F=1$

Truth Table for F

A	B	C	F
0	0	0	0
0	0	1	1 $\rightarrow \bar{A}\bar{B}C \equiv m_1$
0	1	0	1 $\rightarrow \bar{A}B\bar{C} \equiv m_2$
0	1	1	0
1	0	0	1 $\rightarrow A\bar{B}\bar{C} \equiv m_4$
1	0	1	0
1	1	0	0
1	1	1	1 $\rightarrow ABC \equiv m_7$

($ABC = 001 = \text{binary } \#1$)

Implement F using Logic Gates



Product of Sums Express Boolean function as AND of ORs

$$\bar{F} = \underbrace{A\bar{B}\bar{C}}_{m_6} + \underbrace{A\bar{B}C}_{m_5} + \underbrace{\bar{A}BC}_{m_3} + \underbrace{\bar{A}\bar{B}C}_{m_0}$$

Think about
where $\bar{F}=0$

$$F = \bar{\bar{F}} = \overline{A\bar{B}\bar{C} + A\bar{B}C + \bar{A}BC + \bar{A}\bar{B}C}$$

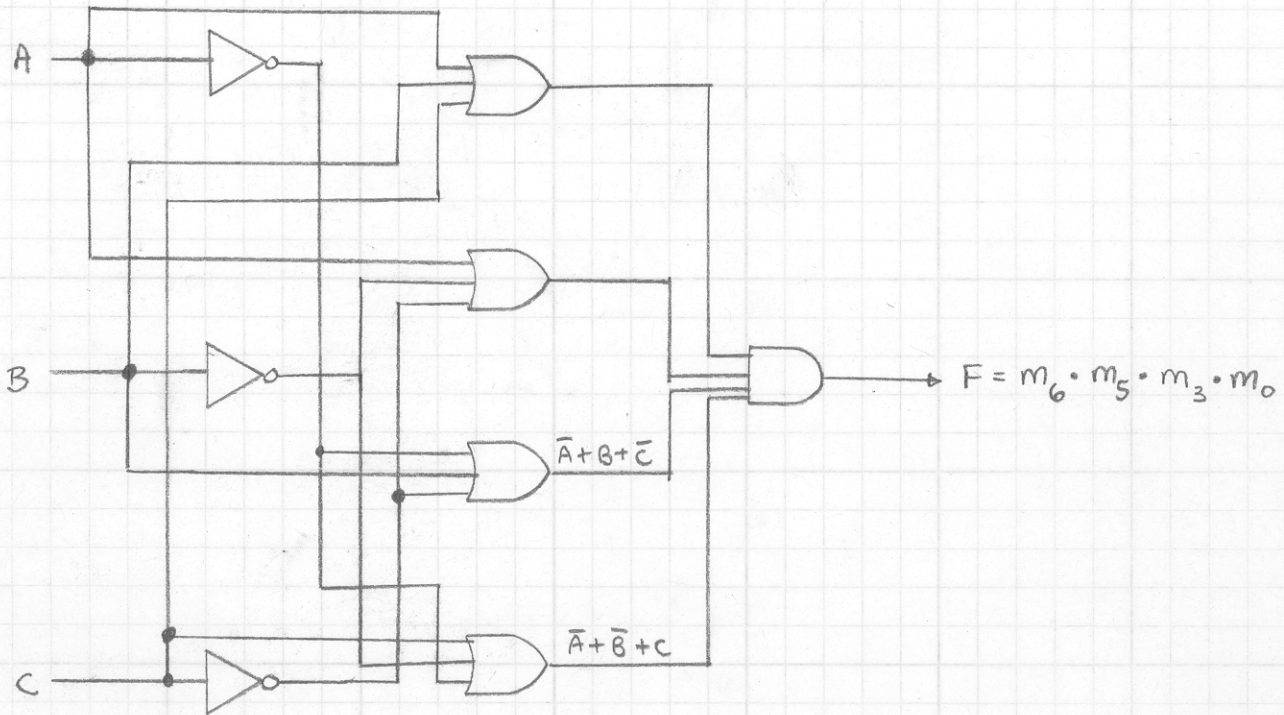
Using De Morgan's Theorem:

$$F = \overline{(A\bar{B}\bar{C})(A\bar{B}C)(\bar{A}BC)(\bar{A}\bar{B}C)} = \underbrace{(\bar{A}+B+C)(\bar{A}+B+\bar{C})(A+\bar{B}+\bar{C})(A+B+C)}_{\text{maxterm}}$$

Truth Table for F Same as Above

A	B	C	F
0	0	0	0 $\rightarrow \bar{A}\bar{B}\bar{C} \equiv m_0$
0	0	1	1
0	1	0	1
0	1	1	0 $\rightarrow \bar{A}BC \equiv m_3$
1	0	0	1
1	0	1	0 $\rightarrow A\bar{B}C \equiv m_5$
1	1	0	0 $\rightarrow AB\bar{C} \equiv m_6$
1	1	1	1

Logic Gate Implementation of F using Product of Sums



III. Karnaugh Maps

Karnaugh maps are a graphical method of simplifying Boolean expressions.

$$\text{Ex: } F = \underbrace{\bar{A}\bar{B}\bar{C}}_{m_0} + \underbrace{\bar{A}\bar{B}C}_{m_1} + \underbrace{\bar{A}B\bar{C}}_{m_2} + \underbrace{AB\bar{C}}_{m_6}$$

First, try to simplify using Boolean algebra:

$$F = \bar{A}\bar{B}(\underbrace{\bar{C}+C}_1) + B\bar{C}(\underbrace{\bar{A}+A}_1) = \bar{A}\bar{B} + B\bar{C}$$

Karnaugh Maps of 3 Variables

F is sum of products (minterms)

A \ BC	00	01	11	10
0	m_0	m_1	m_3	m_2
1	m_4	m_5	m_7	m_6

$$\text{Ex: } F = \bar{A}\bar{B}\bar{C} + \bar{A}\bar{B}C + \bar{A}B\bar{C} + AB\bar{C}$$

A \ BC	00	01	11	10
0	1	1		1
1				1

Annotations: $\bar{A}\bar{B}$ (covering the first two cells of row 0), $\bar{A}\bar{C}$ (covering the first two cells of row 0 and the last cell of row 1), $B\bar{C}$ (covering the last cell of row 0 and the last cell of row 1).

Only need enough terms to "cover" all 1's:

$$F = \bar{A}\bar{B} + B\bar{C}$$

For minterms, think where is $F=1$?

Karnaugh Maps of 4 Variables

AB \ CD	00	01	11	10
00	m_0	m_1	m_3	m_2
01	m_4	m_5	m_7	m_6
11	m_{12}	m_{13}	m_{15}	m_{14}
10	m_8	m_9	m_{11}	m_{10}

Handwritten Karnaugh map for the function $F(A, B, C, D) = \sum(0, 1, 2, 3, 4, 5, 6, 7)$. The map is a 4x4 grid with rows labeled AB (00, 01, 11, 10) and columns labeled CD (00, 01, 11, 10). All cells contain the value 1. Groupings are shown with loops: a vertical loop for A (covering columns 00 and 01), a horizontal loop for B (covering rows 00 and 01), and a horizontal loop for C (covering columns 11 and 10).

AB \ CD	00	01	11	10
00				
01	1	1		
11			1	
10	1	1		

$$F = \overline{A}B\overline{C} + A\overline{B}\overline{C} + ABCD$$

Karnaugh Maps for Maxterms (Product of Sums)

Group "0"s in Karnaugh Map.

$$\begin{aligned}
 \text{Ex: } F_1 &= \underbrace{(A+B+C+D)}_{m_0} \underbrace{(A+B+C+\bar{D})}_{m_1} \underbrace{(A+\bar{B}+C+D)}_{m_4} \cdot \\
 &\quad \underbrace{(A+\bar{B}+\bar{C}+D)}_{m_6} \underbrace{(A+\bar{B}+\bar{C}+\bar{D})}_{m_7} \underbrace{(\bar{A}+B+C+\bar{D})}_{m_9} \cdot \\
 &\quad \underbrace{(\bar{A}+\bar{B}+C+D)}_{m_{12}} \underbrace{(\bar{A}+\bar{B}+\bar{C}+D)}_{m_{14}} \underbrace{(\bar{A}+\bar{B}+\bar{C}+\bar{D})}_{m_{15}}
 \end{aligned}$$

Truth Table for F_1

A	B	C	D	F_1	
0	0	0	0	0	m_0 $A+B+C+D$
0	0	0	1	0	m_1 $A+B+C+\bar{D}$
0	0	1	0	1	m_2
0	0	1	1	1	m_3
0	1	0	0	0	m_4 $A+\bar{B}+C+D$
0	1	0	1	1	m_5
0	1	1	0	0	m_6 $A+\bar{B}+\bar{C}+D$
0	1	1	1	0	m_7 $A+\bar{B}+\bar{C}+\bar{D}$
1	0	0	0	1	m_8
1	0	0	1	0	m_9 $\bar{A}+B+C+\bar{D}$
1	0	1	0	1	m_{10}
1	0	1	1	1	m_{11}
1	1	0	0	0	m_{12} $\bar{A}+\bar{B}+C+D$
1	1	0	1	1	m_{13}
1	1	1	0	0	m_{14} $\bar{A}+\bar{B}+\bar{C}+D$
1	1	1	1	0	m_{15} $\bar{A}+\bar{B}+\bar{C}+\bar{D}$

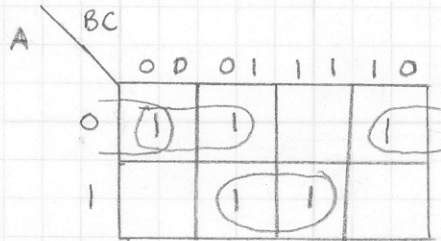
AB \ CD	00	01	11	10
00	0	0		
01	0		0	0
11	0		0	0
10		0		

$$F1 = (\bar{B} + D)(\bar{B} + \bar{C})(B + C + \bar{D})(A + B + C)$$

I. Last Word on Karnaugh Maps

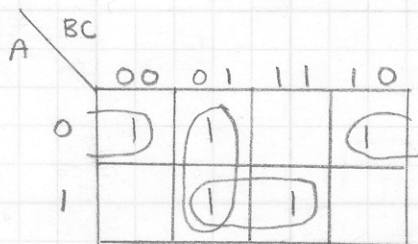
What does it mean to simplify an expression?

Ex: Find simplest expression for following Karnaugh map:



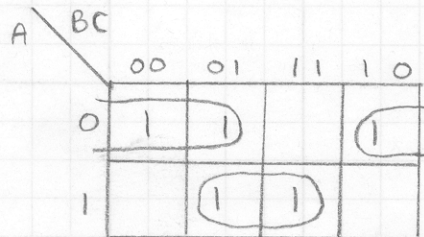
One grouping: $\bar{A}\bar{C} + AC + \bar{A}\bar{B}$

3 terms, 6 "literals" (signals such as A, B, \bar{B} , etc.)

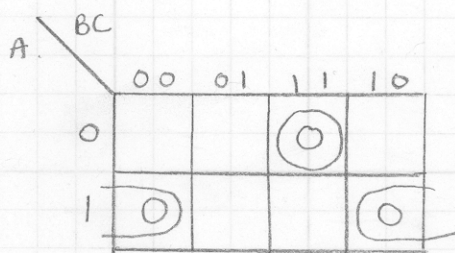


Alternative grouping: $AC + \bar{A}\bar{C} + \bar{B}C$

3 terms, 6 literals



$\bar{A} \cdot (\bar{B} + \bar{C}) + AC$ 5 literals, but no longer sum of products



$(A + \bar{B} + \bar{C}) \cdot (\bar{A} + C)$ 5 literals

Consider circuit implementation: (starting with A, B, C)

$\bar{A}\bar{C} + AC + \bar{A}\bar{B} \rightarrow$ 3 inverters, 3 2-input ANDs, 1 3-input OR

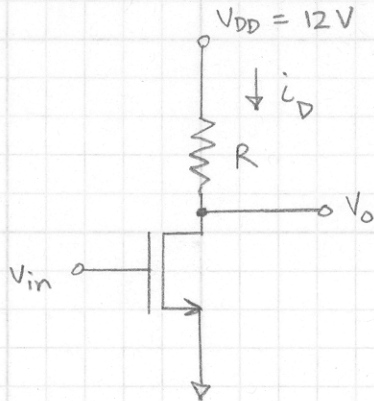
$AC + \bar{A}\bar{C} + \bar{B}C \rightarrow$ 3 inverters, 3 2-input ANDs, 1 3-input OR

$\bar{A} \cdot (\bar{B} + \bar{C}) + AC \rightarrow$ 3 inverters, 2 2-input ANDs, 2 2-input ORs

$(A + \bar{B} + \bar{C}) \cdot (\bar{A} + C) \rightarrow$ 3 inverters, 1 2-input AND, 1 2-input OR, 1 3-input OR

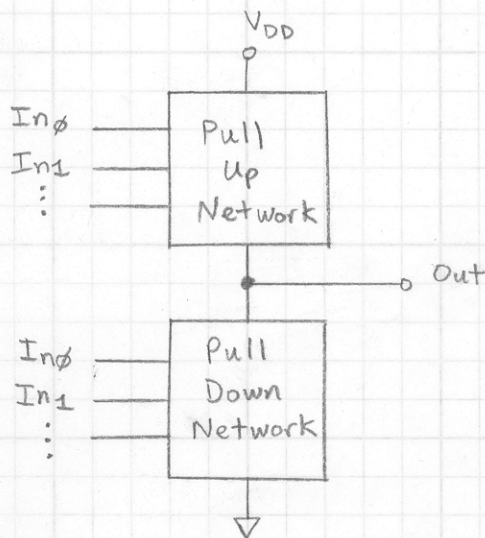
II. MOSFET Logic Gates

NMOS Inverter with Resistor Pull-Up



Truth Table:	V _{in}	V _{out}
	0	1
	1	0

General Logic Circuit Concept



In_0	In_1	\dots	In_N	Out
0	0	\dots	1	0 \leftarrow pull down network
1	0	\dots	1	0
1	1	\dots	1	1 \leftarrow pull up network

Each network can be implemented with transistors, resistors, sometimes diodes, etc.