## SINGLE-BIT MEMORIES (FLIP-FLOPS)

## Semiconductor Memories

- "Single-bit" Memory (Foreground)
- Individual memory circuits that store a single bit of information and have at least the following I/O ports:

1) data input and 2) data output
1. "Clock-less" latch
2. Transparent Level-Sensitive Latch
3. Edge-triggered Flip-Flop

- "Array" Memory (Background)
- Large memory circuit that stores many bits of information
 organized into multiple words accessed by an address


1. SRAM
2. Multi-ported SRAM
3. DRAM
4. Flash
5. etc.

## Instantiating Flip-Flops/Registers

- One way to build a FF/register (do not use this)

```
reg a;
always @(posedge clk) begin
    a = a_c;
end
```

- The " $=$ " is a "blocking assignment" which
 causes the simulator to "block" on an assignment until the operation is completed, then it moves to the next statement
- It makes a race condition possible

```
reg b, c;
always @(posedge clk) begin
    b = a;
    c = b;
end
```



- In this case, $a$ races to $c$ in one cycle!


## Instantiating Flip-Flops/Registers

- One solution would be to simply reorder the reg assignment commands to eliminate the race
- $\quad a$ races to $c$ in one cycle (broken)

```
reg b, c;
always @(posedge clk) begin
            b = a;
            c = b;
end
```

- $\quad a$ progresses to $b$ and $b$ to $c$ in one cycle each as expected (ok)

```
reg b, c;
always @(posedge clk) begin
    c = b; // reordered from above
    b = a; // reordered from above
end
```

- Of course this is a terrible solution and you should never use it. Besides being very tedious and prone to errors, it will not work with loops.



## Instantiating Flip-Flops/Registers

- The correct solution is to use a "non-blocking assignment" written with "<=" which causes the simulator to evaluate the right side of the expression when the statement is encountered, but the assignment of the left side is not done until the end of that time step in "simulator time"
- With the verilog below, the registers perform as normal FFs behave without a race regardless of their ordering

```
reg b, c;
always @(posedge clk) begin
    b <= a;
    c}<= b
end
```

The key differentiator is the time when the value $b$ is sampled. Here it is not "blocked" by the previous assignment $b<=a$


## Single-Bit Memory Rule \#1 (out of 9 total rules)

- Rule \#1 (always follow in this class):

Use only edge-triggered flip-flops and never transparent (clock level-sensitive) latches

- Edge-triggered flip-flops are generally robust with regards to clocking
- Transparent latches are vulnerable to signals racing through more than one latch during a
 single clock pulse; solutions are generally non-trivial
- Requiring a minimum delay between latches works but adds area and power dissipation
- Requiring a maximum pulse width on clocks (likely $\ll 50 \%$ ) can be problematic
- A robust but tedious solution is to use two clocks with different nonoverlapping phases, and require consecutive latches to use clocks of differing phases




## Single-Bit Memory Rule \#2

- Rule \#2 (always follow in this class):

For combinational logic always blocks, always use blocking assignments ("=")

```
// OR gate
always @(a or b) begin
    c = a | b;
end
```


## Single-Bit Memory Rule \#3

- Rule \#3 (always follow in this class):

For flip-flop (register) always blocks, always use non-blocking assignments ("<=")

```
always @(posedge clk) begin
    sum <= #1 c_sum;
    r_product <= #1 product;
end
```


## Single-Bit Memory Rule \#4

- Rule \#4 (always follow in this class):

Add "\#1" to give one unit of clock-to-Q delay to increase waveform readability

- Synthesis tools have no way to implement a \#1 delay so they typically give a warning saying they can not add the delay and that the tool has ignored it
- This warning can be ignored (the only warning that can be automatically ignored!)
- Interestingly, Quartus does not give a warning



## Single-Bit Memory Rule \#5

- Rule \#5 (really a guideline): Normally do not include any logic in flip-flop declarations. You will be tempted to include functions such as resets, initializations, and counter incrementing. Try to resist but it is ok if logic is simple.

```
// Flip-flop declarations
always @(posedge clock) begin
    state <= #1 state_c;
    count <= #1 count_c;
    if (reset == 1'b1 && state_pipe3 != 8'hB2) begin
        state <= #1 4'b0000; // init state
    end
end
```


## Single-Bit Memory Rule \#6

- Rule \#6 (really a suggestion):
- For large designs, it is usually clearer to group related combinational logic and FFs in separate areas of the module

```
// Block 1
[combinational logic]
always @(posedge clock) begin
    count <= #1 count_c;
end
// Block 2
[combinational logic]
always @(posedge clock) begin
    data_r <= #1 data;
end
// Block 3
[combinational logic]
always @(posedge clock) begin
    cat <= #1 mouse;
end
```


## Preview of 2 Clock-Related Rules

Rule \#7

- These will be covered in more detail in the Clock section

1) Only clock signals may connect to flip-flop or latch clock inputs
2) Clock signals may not connect to any node other than a flip-flop or latch clock input

- No logic gate inputs
- No flip-flop or latch inputs other than the clock input

3) There are only a few exceptions

## Rule \#8: Signal Naming Conventions

- It is helpful to have conventions for signal names
- Easier for others to understand your code
- Easier for YOU to understand your code
- Add a suffix to signal names to indicate they are from an earlier pipeline stage or a later pipeline stage
- *_c - input to a register (e.g., sum_c)
- *_r-output of a register (e.g., sum_r)
(combinational logic version of sum)



## Signal Naming Conventions

- Another possibility I have seen is to add a prefix but this has the possibly-negative feature that associated signals are not adjacent when sorted alphabetically
- $c_{-}^{*}$ - input to a register (e.g., $c_{-}$sum)
- $r_{-}{ }^{*}$ - output of a register (e.g., $r_{-}$sum)



## Signal Naming Conventions

- If a signal is pipelined across multiple pipe stages, it is probably a good idea to indicate that in the signal's name-for example with a suffix such as _pipeX
- inA_pipe1 - inA signal in pipeline stage 1



## Reset-able and Enable-able Registers

- Sometimes it is convenient or necessary to have flip-flops with special inputs like reset and enable
- When designing flip-flops/registers, it is ok (possibly required) for there to be cases where the always block is entered, but the reg is not assigned-such as when a FF is disabled
- No fancy code, just make it work
- Normally use synchronous reset instead of asynchronous reset (Rule \#9 later)



## Reset-able Registers

- Two example FFs with active-high reset
- The code also leaves out begin-end statements which is not ideal
- This code "bends" Guideline \#5

```
always @(posedge clk) begin
    if (reset == 1'b0)
        out <= #1 D;
    else
        out <= #1 1'b0;
end
always @(posedge clk) begin
    if (reset == 1'b1)
        out <= #1 1'b0;
    else
        out <= #1 D;
end
```


## Reset-able Registers

- Example FF registers with reset
- This code bends Guideline \#5

```
reg out;
reg [7:0] count;
reg [23:0] rgb;
always @(posedge clk) begin
    out <= #1 D;
    count <= #1 count_c;
    rgb <= #1 color24;
    if (reset == 1'b1) begin
        out <= #1 1'b0;
        count <= #1 8'b0000_0000; // ideally put reset elsewhere
        // assume rgb does not need to be reset
    end
end
```


## Preset-able Registers

- Example FF with preset
- This code bends Guideline \#5

```
always @(posedge clk) begin
    out <= #1 D;
    if (preset == 1'b1) begin
        out <= #1 1'b1;
    end
end
```


## Enable-able Registers

- Example FF with reset
- This is not combinational logic so it is not necessary that all outputs are set in every possible path through the always block
- This code bends Guideline \#5

```
always @(posedge clk) begin
    if (enable == 1'b1) begin
        out <= #1 D;
    end
end
```

- Compare with Common Mistake \#4


## Reset-able and Enable-able Registers

- Example FF with reset and enable (reset has priority)

| reset | enable | Action |
| :---: | :---: | :--- |
| 0 | 0 | Do nothing |
| 0 | 1 | D Flip-Flop |
| 1 | 0 | Reset? or Do nothing? |
| 1 | 1 | Reset, Q=0 |

always @(posedge clk) begin
if (reset == 1'b1) // highest priority out $<=$ \#1 1'b0;
else if (enable == 1'b1) out <= \#1 c_out;
// ok if no assignment (out holds value)
end

## Reset-able and Enable-able Registers

- Example FF with reset and enable (enable has priority)

```
\begin{tabular}{cc|l} 
reset & enable & Action \\
\hline 0 & 0 & Do nothing \\
0 & 1 & D Flip-Flop \\
1 & 0 & Reset? or Do nothing? \\
1 & 1 & Reset, Q=0
\end{tabular}
always @(posedge clk) begin
    if (enable == 1'b1) begin // highest priority
        if (reset == 1'b1)
            out <= #1 1'b0;
        else
            out <= #1 c_out;
    end
    // It is ok if there is no assignment to "out" in some
    // cases (out is not assigned when enable==0 in which
    // case out holds its value). Recall that this is never
    // ok for combinational logic.
end
```


## Resets Within a Large Datapath

- Use reset-able FFs only where truly needed
- Reset-able FFs are a little larger and higher power
- Requires the global routing of the high-fanout reset signal



## Asynchronous Resets

- Asynchronous resets will react to any pulse, even a very short "glitch" in the middle of a clock period
- What if the reset signal comes directly from combinational logic where glitches are common?
- What if a glitch comes from capacitive coupling or inductive coupling of wires?
- How to test in a big chip?
- They utilize a different circuit than what a synchronous reset uses
- However they are needed in:
- clock generation
- asynchronous interface logic
- Rule \#9: Use only synchronous resets (always follow in this class)


